

Univalent Foundations of Mathematics and Homotopical Algebra

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Homotopical Algebra and Geometry
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April 10th, 2015

The univalent foundations of mathematics programme

Origin:

- ▶ formulated around 2009 by Vladimir Voevodsky
- ▶ related to Homotopy Type Theory

Motivation:

- ▶ to facilitate computer-assisted verification of proofs

Main feature: it combines ideas from

- ▶ type theory
- ▶ homotopy theory

Overview of the talk

Part I: Simplicial sets

- ▶ simplicial sets
- ▶ univalent fibrations
- ▶ simplicial univalence

Part II: Univalent foundations

- ▶ type theory
- ▶ the univalence axiom
- ▶ univalent foundations

Part I: Simplicial sets

Simplicial sets as a category

SSet = category of simplicial sets.

SSet is a presheaf topos \Rightarrow

- ▶ all small limits and colimits exist
- ▶ it is locally cartesian closed: all slices are cartesian closed. Equivalently:

$$\begin{array}{c} X \\ \downarrow f \\ Y \end{array}$$

$$\begin{array}{ccc} & \mathbf{SSet}/X & \\ \Sigma_f \swarrow \lrcorner & \uparrow \Delta_f & \lrcorner \searrow \Pi_f \\ & \mathbf{SSet}/Y & \end{array}$$

Simplicial sets as a model category

SSet has a model structure $(\mathcal{W}, \mathcal{F}, \mathcal{C})$, where

- ▶ \mathcal{W} = weak homotopy equivalences
- ▶ \mathcal{F} = Kan fibrations
- ▶ \mathcal{C} = monomorphisms

In particular, every diagram

$$\begin{array}{ccc} X & \longrightarrow & E \\ \downarrow i & \nearrow & \downarrow p \\ Y & \longrightarrow & B \end{array} \quad \text{where } \begin{cases} i \in \mathcal{W} \cap \mathcal{C} \\ p \in \mathcal{F} \end{cases}$$

has a diagonal filler.

The fibrant objects are exactly the Kan complexes.

The model structure is cofibrantly generated and (left and right) proper:

- ▶ right proper = pullback along fibrations preserves weak equivalences
- ▶ left proper = pushout along cofibrations preserves weak equivalences

Homotopy theory in simplicial sets

The rich structure of **SSet** allows us to internalize a lot of constructions.

Example. Let $p: E \rightarrow B$ a fibration. There is a fibration

$$(s, t): \text{Weq}(E) \rightarrow B \times B$$

such that the fiber over (x, y) is

$$\text{Weq}(E)_{x,y} = \{w: E_x \rightarrow E_y \mid w \in \mathcal{W}\}$$

Note. Given $p: E \rightarrow B$, we have

$$\begin{array}{ccc} B & \xrightarrow{\quad} & \text{Weq}(E) \\ i \downarrow & \nearrow j_p & \downarrow (s,t) \\ \text{Path}(B) & \xrightarrow{\quad} & B \times B \end{array}$$

Univalent fibrations

Definition. A fibration $p: E \rightarrow B$ is said to be **univalent** if

$$j_p: \text{Path}(B) \rightarrow \text{Weq}(E)$$

is a weak equivalence.

Idea. Weak equivalences between fibers are ‘witnessed’ by paths in the base.

Proposition (Kapulkin, Lumsdaine and Voevodsky). A fibration $p: E \rightarrow B$ is univalent if and only if for every fibration $p': E' \rightarrow B'$, the space of squares

$$\begin{array}{ccc} E' & \xrightarrow{u} & E \\ p' \downarrow & & \downarrow p \\ B' & \xrightarrow{v} & B \end{array}$$

such that

$$(u, p'): E' \rightarrow B' \times_B E$$

is a weak equivalence, is either empty or contractible.

Idea. Essential uniqueness of u, v (when they exist).

The generic Kan fibration

Fix an inaccessible cardinal κ . There exists a fibration

$$\pi: \tilde{U} \rightarrow U$$

that weakly classifies fibrations with fibers of cardinality $< \kappa$, i.e. for every such $p: E \rightarrow B$ there exists a pullback diagram

$$\begin{array}{ccc} E & \longrightarrow & \tilde{U} \\ p \downarrow & & \downarrow \pi \\ B & \longrightarrow & U. \end{array}$$

Note. Given $x: 1 \rightarrow U$, we can form a pullback

$$\begin{array}{ccc} \text{El}(x) & \longrightarrow & \tilde{U} \\ \downarrow & & \downarrow \\ 1 & \xrightarrow{x} & U \end{array}$$

We think of x as the 'name' of the Kan complex $\text{El}(x)$.

Simplicial univalence

Theorem (Voevodsky). The generic fibration $\pi : \tilde{U} \rightarrow U$ is univalent.

Several proofs:

- ▶ Voevodsky
- ▶ Lumsdaine, Kapulkin, Voevodsky (using simplifications by Joyal)
- ▶ Moerdijk (using fiber bundles)
- ▶ Cisinski (general setting)

Note. The fibration $\pi : \tilde{U} \rightarrow U$ is therefore

versal & univalent

for the class of κ -small fibrations.

The universe is Kan

Theorem. The codomain of $\pi : \tilde{U} \rightarrow U$ is a Kan complex.

Proof. Show

$$\begin{array}{ccc} \Lambda_n^k & \xrightarrow{\forall f} & U \\ i \downarrow & & \nearrow \exists f' \\ \Delta[n] & & \end{array}$$

This reduces to the problem of extending fibrations along horn inclusions:

$$\begin{array}{ccc} E & \cdots \cdots \cdots \rightarrow & ? \\ p \downarrow & & \downarrow \\ \Lambda_n^k & \xrightarrow{i} & \Delta[n] \end{array}$$

which can be done using the theory of minimal fibrations.

Note. A similar unfolding is possible for the univalence of $\pi : \tilde{U} \rightarrow U$.

Part II: Univalent Foundations

General idea

Type theories are formal systems.

They have axioms for manipulating types and their elements:

$$A \text{ type, } a \in A.$$

Many type theories in the literature.

Key fact. The axioms of Martin-Löf type theories correspond quite closely to a fragment of the structure of **SSet**.

Original inspiration for Martin-Löf type theories comes from proof theory and theoretical computer science (cf. implementation in Coq, Agda).

We have axioms for

$$0, 1, \dots, \mathbb{N}$$

$$X \times Y, X + Y, Y^X, \dots$$

$$\text{Id}_X(x, y), \sum_{b \in B} E(b), \prod_{b \in B} E(b), \cup$$

A dictionary

| Simplicial sets | Type Theory |
|--|------------------------------------|
| Kan complexes | Types |
| $x: 1 \rightarrow X$ | $x \in X$ |
| p is a path from x to y in X | $p \in \text{Id}_X(x, y)$ |
| fibration $p: E \rightarrow B$ | $b \in B \vdash E(b)$ type |
| total space of $p: E \rightarrow B$ | $\sum_{b \in B} E(b)$ |
| space of sections of $p: E \rightarrow B$ | $\prod_{b \in B} E(b)$ |
| base of the generic fibration $\pi: \tilde{U} \rightarrow U$ | type universe U |
| the generic fibration $\pi: \tilde{U} \rightarrow U$ | $x \in U \vdash \text{El}(x)$ type |

Note. We do not have counterparts for all the structure of simplicial sets. This is a limitation, but ensures good proof-theoretical and computational properties.

Type theory and homotopical algebra (I)

Theorem (Awodey and Warren). The axioms for identity types can be stated equivalently as follows:

1. For every type X and $x, y \in X$, we have a type

$$\text{Id}_X(x, y).$$

2. For every X , the diagonal map $\Delta_X : X \rightarrow X \times X$ has a factorisation

$$X \xrightarrow{i} \sum_{x, y \in X} \text{Id}_X(x, y) \xrightarrow{p} X \times X.$$

3. Every commutative diagram of the form

$$\begin{array}{ccc} X & \longrightarrow & \sum_{b \in B} E(b) \\ \downarrow i & & \downarrow \pi_1 \\ \sum_{x, y \in X} \text{Id}_X(x, y) & \longrightarrow & B. \end{array}$$

has a diagonal filler.

Type theory and homotopical algebra (II)

Theorem (Gambino and Garner). There is a weak factorization system $(\mathcal{L}, \mathcal{R})$, where

$$\mathcal{L} = \text{functions with LLP w.r.t. every } \pi_1 : \sum_{b \in B} E(b) \rightarrow B.$$

Theorem (Voevodsky). The category **SSet** is a model of Martin-Löf type theory and of the Univalence Axiom.

Theorem (Garner and van den Berg; Lumsdaine). Every type X determines a weak ∞ -groupoid $\Pi_\infty(X)$ (in the sense of Batanin-Leinster) in which

- ▶ objects are elements of X ,
- ▶ 1-cells $p : x \rightarrow y$ are $p \in \text{Id}_X(x, y)$,
- ▶ 2-cells $\alpha : p \rightarrow q$ are $\alpha \in \text{Id}_{\text{Id}_X(x, y)}(p, q)$,
- ▶ ...

Univalent foundations of mathematics

These connections suggest to:

1. use type theory as a language for speaking about homotopy types,
2. develop mathematics using this language; in particular

$\text{sets} =_{\text{def}} \text{discrete homotopy types},$

3. add axioms to type theory motivated by homotopy theory

Homotopy-theoretic notions in type theory

- ▶ A type X is said to be **contractible** if the type

$$\text{iscontr}(X) =_{\text{def}} \sum_{x \in X} \prod_{y \in X} \text{Id}_X(x, y)$$

is inhabited.

- ▶ The **homotopy fiber** of $f : X \rightarrow Y$ at $y \in B$ is the type

$$\text{hfiber}(f, y) =_{\text{def}} \sum_{x \in X} \text{Id}_Y(fx, y).$$

- ▶ We say that a function $f : X \rightarrow Y$ is a **weak equivalence** if for every $y \in Y$, the homotopy fiber of f at y is contractible.
- ▶ For types X and Y , there is a type

$$\text{Weq}(X, Y)$$

of weak equivalences from X to Y .

The hierarchy of homotopy levels in type theory

Definition. We say that a type X has

- ▶ **homotopy level 0** if it is contractible,
- ▶ **homotopy level $n + 1$** if $\text{Id}_X(x, y)$ is an n -type for all $x, y \in X$.

Example. Let X be a type.

X has level 1 \Leftrightarrow for all $x, y \in X$, $\text{Id}_X(x, y)$ is contractible
 \Leftrightarrow if X is inhabited, then it is contractible.

Idea:

| | | | | |
|--------------------|---|--------------------|-----------|-------------------|
| <i>Level</i> | 0 | 1 | 2 | 3 |
| <i>Types</i> | * | $\emptyset, \{*\}$ | sets | groupoids |
| <i>Mathematics</i> | | “logic” | “algebra” | “category theory” |

The univalence axiom

The type universe U is such that if $x \in U$ then $\text{El}(x)$ is a type.

Idea.

$$\begin{array}{ccc} \text{El}(x) & \longrightarrow & \sum_{x \in U} \text{El}(x) \\ \downarrow & & \downarrow \pi_1 \\ \mathbf{1} & \xrightarrow{x} & U \end{array}$$

Note. For $x, y \in U$, we have

- ▶ the type of paths from x to y , $\text{Id}_U(x, y)$
- ▶ the type of weak equivalences from $\text{El}(x)$ to $\text{El}(y)$, $\text{Weq}(\text{El}(x), \text{El}(y))$

and there is a canonical map

$$j_{x,y} : \text{Id}_U(x, y) \rightarrow \text{Weq}(\text{El}(x), \text{El}(y)).$$

Univalence Axiom. For all $x, y \in U$, the map $j_{x,y}$ is a weak equivalence.

Remarks on the univalence axiom

The univalence axiom is valid in **SSet** by the univalence of $\pi : \tilde{U} \rightarrow U$.

This axiom has several interesting aspects from a logical point of view:

- ▶ it forces the type universe U not to be of level 2 (the level of “sets”)
- ▶ it is not valid in the set-theoretical model
- ▶ it allows us to treat isomorphic structures as if they were equal
- ▶ it has useful consequences, such as function extensionality (Voevodsky)
- ▶ it gives the ‘Rezk completion’ of a category (Ahrens, Kapulkin, Shulman)

Further aspects

- ▶ Synthetic homotopy theory (Shulman, Lumsdaine, Licata, Brunerie, ...)
- ▶ Higher inductive types (Shulman, Lumsdaine)
- ▶ Homotopy-initial algebras in type theory (Awodey, Gambino and Sojakova)
- ▶ Other models of univalent foundations (Coquand et al., Shulman, Cisinski)
- ▶ Univalent maps in $(\infty, 1)$ -categories (Gepner and Kock)