

Forking in accessible categories

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joint work with M.Lieberman and S. Vasey

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This is open even for $\mu = \aleph_0$. Is the category of complete metric spaces and isometries a counterexample?

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Theorem 1. Locally multipresentable categories whose morphisms are monomorphisms coincide with universal μ -AECs.

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A *polyinitial* object is a set \mathcal{I} of objects of a category \mathcal{K} such that for every object M in \mathcal{K} :

1. There is a unique $i \in \mathcal{I}$ having a morphism $i \rightarrow M$.
2. For each $i \in \mathcal{I}$, given $f, g : i \rightarrow M$, there is a unique (isomorphism) $h : i \rightarrow i$ with $fh = g$.

We get groups of automorphisms of members of a polyinitial object. In the case of a multinitial objects, they are singletons. An example of a locally polypresentable category whose morphisms are monomorphisms are algebraically closed fields. The polyinitial object is formed by algebraic closures of the multiinitial object in fields.

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Lemma 1. Let \mathcal{K} be a coregular locally μ -presentable category and \mathcal{K}_{reg} be the category having the same objects as \mathcal{K} and regular monomorphisms as morphisms. Then \mathcal{K}_{reg} is locally μ -multipresentable.

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Examples of coregular locally presentable categories: Grothendieck toposes, Grothendieck abelian categories, **Gra** graphs, **Gr** groups, **Bool** Boolean algebras, **Ban** Banach spaces with linear contractions, **Hilb** Hilbert spaces with linear isometries, **CAlg** commutative unital C^* -algebras, etc.

Let \mathcal{K} be a coregular locally presentable category and

$$\begin{array}{c} M_1 \\ \uparrow \\ M_0 \rightarrow M_2 \end{array}$$

a span in \mathcal{K}_{reg} . Let

$$\begin{array}{ccc} M_1 & \longrightarrow & P \\ \uparrow & & \uparrow \\ M_0 & \longrightarrow & M_2 \end{array}$$

a pushout in \mathcal{K} . Then a multipushout in \mathcal{K}_{reg} is formed by all squares

$$\begin{array}{ccc} M_1 & \longrightarrow & Q \\ \uparrow & & \uparrow \\ M_0 & \longrightarrow & M_2 \end{array}$$

where the induced morphism $P \rightarrow Q$ is an epimorphism.

The *notion of independence* \perp in \mathcal{K} consists in the choice of squares

$$\begin{array}{ccc} M_1 & \rightarrow & M_3 \\ \uparrow & & \uparrow \\ M_0 & \rightarrow & M_2 \end{array}$$

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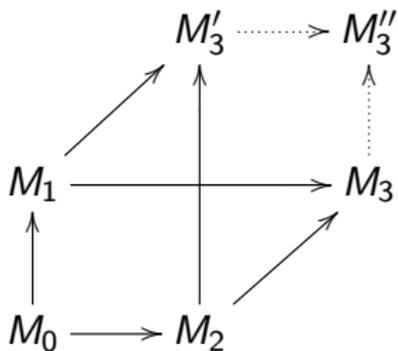
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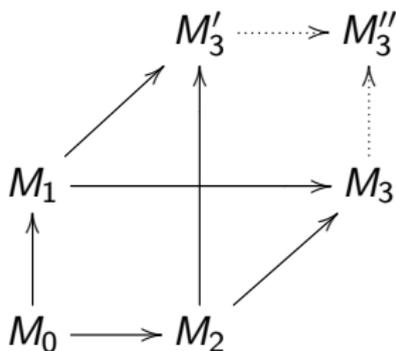
The following properties should be satisfied

- (i) invariance under isomorphisms of squares
- (ii) independence on M_3 ,
- (iii) existence,
- (iv) uniqueness,
- (v) symmetry,
- (vi) closedness under compositions of squares, and
- (vii) accessibility.

(ii) means to be closed under the equivalence generated by

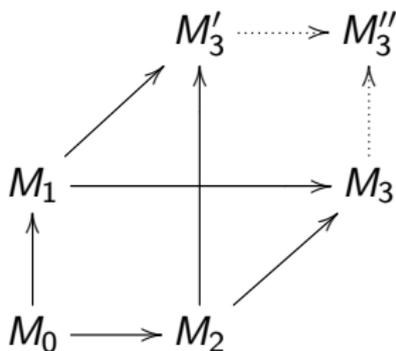


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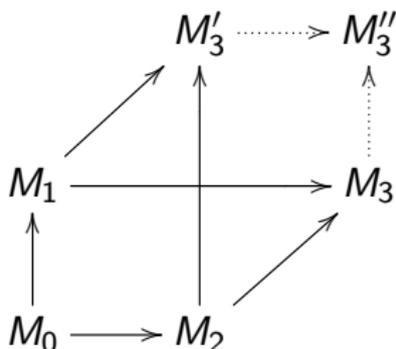
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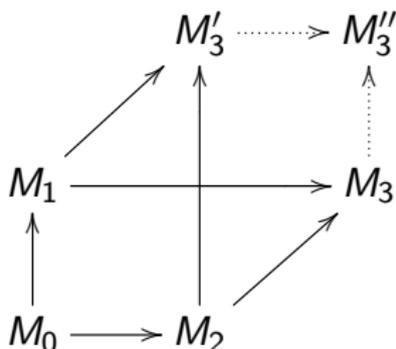


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(vii) the category whose objects are morphisms in \mathcal{K} and whose morphisms are independent squares is accessible.

Theorem 2. Let \mathcal{K} be an accessible category with chain bounds whose morphisms are monomorphisms. Then \mathcal{K} has at most one notion of independence.

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Theorem 3. Let \mathcal{K} be an accessible category whose morphisms are monomorphism having a notion of independence. Then \mathcal{K} is tame, stable and does not have the order property.

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Theorem 4. Let \mathcal{K} be a coregular locally presentable category with effective unions. Then \mathcal{K}_{reg} has an independence notion (consisting of pullback squares).

\mathcal{K} has effective unions if whenever we have a pullback

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and a pushout

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Gra, **Gr**, **Ban**, **Bool** or **CAlg** do not have effective unions. They do not have a notion of independence because they have the order property.

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M_3 does not contain any cross-edge between M_1 and M_2 . Thus we do not have enough independent squares with M_0 and M_2 small. Thus **Gra** does not have an independence notion.

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There is another attempt of independence where we include all cross-edges between M_1 and M_2 . This yields \perp satisfying (i-vi). By Theorem 2, **Gra** does not have an independence notion.

The category \mathcal{K} of graphs whose vertices have order $\leq n$ form a locally finitely presentable category which does not have effective unions. But effective pullback squares provide an independence in \mathcal{K}_{reg} . The reason is that we can add all cross-edges to M_0 and M_2 and keeping them finite.

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Let \mathcal{K} be the category of locally finite graphs, i.e., graphs such any vertex has a finite degree. This category is coregular and locally \aleph_1 -multipresentable. Again, effective pullback squares form a notion of independence in \mathcal{K}_{reg} .

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\mathcal{K} does not have chain bounds but the proof of Theorem 2 still goes through – thus this independence is unique.

Based on Malliaris and Shelah 2011, we say that a graph is stable if it does not contain a copy of the half graph (the bipartite graph on $\mathbb{N} \times \mathbb{N}$ such that $E(i, j)$ iff $i < j$). The category \mathcal{K} of stable graphs is locally \aleph_1 -multipresentable and effective pullback squares do not form an independence notion in \mathcal{K}_{reg} . But we do not know whether \mathcal{K}_{reg} has an independence notion. \mathcal{K} does not have chain bounds and we could not adapt the proof of Theorem 2 to this case.

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Let \mathcal{K} be a locally polypresentable category whose morphisms are monomorphisms having a stable independence notion. Then, for each span, exactly one instance of a polypushout is independent. Moreover, a morphism of spans

$$(\text{id}_{M_0}, h_1, h_2) : (M_0, M_1, M_2) \rightarrow (M_0, M'_1, M'_2)$$

induces a morphism of independent instances of polypushouts. Thus the independence yields a coherent choice of polypushouts.

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Corollary 2. Let \mathcal{K} be a locally polypresentable category with the amalgamation property, chain bounds and whose morphisms are monomorphisms. If \mathcal{K} has two distinct coherent choices of polypushouts then it does not have a notion of independence.

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Consequently, **Gr**, **Ban**, **Bool** and **CAlg** do not have a notion of independence. **Gr** do not have enough regular injectives and thus regular monomorphisms cannot be cofibrantly generated. In all other cases, regular injectives do not form an accessible category and thus regular monomorphisms cannot be cofibrantly generated again.

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Another consequence is that regular monomorphisms in **Gra** are not cofibrantly generated. Equivalently, **Gra** does not have enough regular injectives or those are not accessible.

Theorem 6. Let \mathcal{K} be an accessible category whose morphisms are monomorphisms having the amalgamation property and chain bounds. Let κ be a strongly compact cardinal. If \mathcal{K} does not have the order property then the full subcategory of \mathcal{K} consisting of κ -saturated objects has an independence notion.

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Theorem 7. Let \mathcal{K} be an accessible category whose morphisms are monomorphisms having an independence notion. Then there exists a regular cardinal κ such that any independent square with M_0 κ -saturated is a pullback square.