

The lattice of algebraically closed sets

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Independence relations

(1/9)

Independence relation

= invariant relation satisfying:

(symmetry)

$$A \perp_C B \iff B \perp_C A$$

(monotonicity/finite character)

$$A \perp_C B \iff A_0 \perp_C B_0 \\ \forall \text{ finite } A_0 \subseteq A, B_0 \subseteq B$$

(full transitivity)

$$A \perp_D B \iff A \perp_D C \text{ and } A \perp_C B, \\ \text{provided that } D \subseteq C \subseteq B$$

(extension)

$$A \perp_D C, C \subseteq B \Rightarrow \exists A' \equiv_{C \cup D} A : A' \perp_D B$$

(local character)

$$\exists \kappa \quad \forall B, \text{ finite } A_0 \quad \exists C \subseteq B: \quad A_0 \perp_C B \\ \text{and } |C| < \kappa.$$

Not required: boundedness / stationarity over models; independence theorem / amalgamation property / chain condition.

Strict independence relations (2/9)

Easy consequence of the axioms:

$$A \perp_C B \iff \text{acl}(A \cup C) \perp_{\text{acl} C} \text{acl}(B \cup C)$$

\perp is strict

$$\iff a \perp_C a \text{ implies } a \in \text{acl} C$$

$$\iff \text{acl}(A \cup C) \cap \text{acl}(B \cup C) = \text{acl} C \\ \text{whenever } A \perp_C B.$$

Acl = lattice of algebraically closed sets

The axioms can be stated in terms of Acl and its automorphisms.

Examples:

- forking (Shelah, Kim)
- o-minimal independence (Pillay, Steinhorn)
- \mathfrak{p} -forking (Scanlon, Onshuus)

Objects of this talk

(3/9)

1. Re-invent the wheel:
Find a strict independence relation that is the correct one for vector spaces and algebraically closed fields (ACF).
2. Name-dropping:
Mention 4 important mathematicians who were involved in the short boom of lattice theory in the 1930s.
3. Understand thorn-forking.

Modularity

(4/9)

Definition (J. v. Neumann)

$$A \perp_C^a B \iff \text{acl}(A \cup C) \cap \text{acl}(B \cup C) = \text{acl} C.$$

Theorem

\perp^a is a strict independence relation
 \iff Acl is modular.

Proof sketch

\perp^a is almost an independence relation and certainly strict.

Only problem:

$$A \perp_D^a B \text{ and } D \subseteq C \subseteq B \implies A \perp_C^a B.$$

This holds \iff Acl is modular.

This works for 1-based theories such as vector spaces. But not for ACF: not modular.

Generalising modularity

(5/9)

An active field in the 1930s.

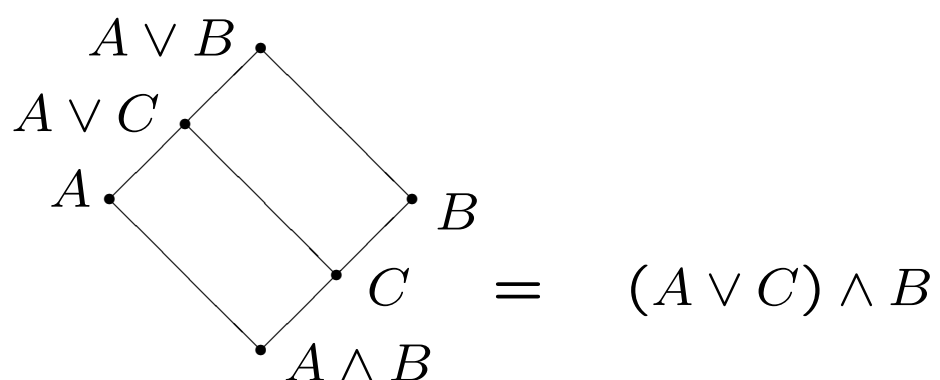
Several attempts:

- Birkhoff's condition (Birkhoff 1935)
- semimodularity (Birkhoff 1935)
- Mac Lane's condition (Mac Lane 1938)
- M-symmetry (Wilcox 1938)

All conditions agree for finite lattices.

All conditions differ in general.

$M(A, B)$ if A and B form a modular pair:



M-symmetry: $M(A, B) \iff M(B, A)$.

M-dividing

(6/9)

Definition (Wilcox 1938)

$$A \perp B \iff A \wedge B = 0 \text{ and } M(A, B).$$

Definition (Scheuermann 1996)

$$A \perp_C^M B \iff \text{acl}(A \cup C) \cap \text{acl}(B \cup C) = \text{acl } C \\ \text{and } M(\text{acl}(A \cup C), \text{acl}(B \cup C)).$$

Theorem (Adler 2005)

\perp^M is a strict independence relation
 \iff Acl is M-symmetric.

Proof idea

Byunghan Kim: symmetry of dividing implies simplicity. Generalise his proof.

M-forking = p-forking

(7/9)

\perp^M is the correct notion for 1-based theories (like vector spaces) and for ACF.

But not preserved under taking reducts!
Saharon Shelah:

A/C does not fork over $D \iff$
 $\forall B \supseteq C \exists A' \equiv_{C \cup D} A :$
 A'/B does not divide over D .

Define M-forking \perp^* from M-dividing \perp^M :

$A \perp_D^* C \iff$
 $\forall B \supseteq C \exists A' \equiv_{C \cup D} A : A' \perp_D^M B.$

Theorem

M-forking = p-forking: $\perp^* = \perp^p$.

Some consequences

(8/9)

Acl is modular $\iff T$ is 1-based rosy
 $\iff \perp^b = \perp^a$.

Acl is M-symmetric
 $\iff T$ rosy and $\perp^b = \perp^M$.

T is rosy
 $\iff \exists$ a strict independence relation
 $\Rightarrow \perp^b$ is **the** coarsest.

Hence T simple $\Rightarrow \perp^b$ is coarser than
forking independence, $A \perp_C^f B \Rightarrow A \perp_C^b B$.

Theorem (Scheuermann 1996)

A strict independence relation with (weak)
canonical bases is **the** coarsest for its theory.

Hence: Forking = \perp -forking
for simple theories with EHI.
(first proof by Clifton Ealy)

Related results

(9/9)

T trivial 1-based rosy \iff Acl is distributive.

T superrosy \iff

T is finitely coded and Acl is arithmetic.

Definitions:

T finitely coded \iff

T rosy and every **global** type is free over a finite set. (Much weaker than superrosy!)

Acl arithmetic

\iff compact elements form a sublattice

\iff they form an ideal

\iff every algebraically closed subset of a finitely generated set is finitely generated.